Automated Intrusion Response CDIS Spring Conference

Kim Hammar

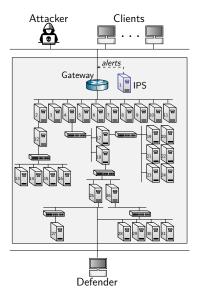
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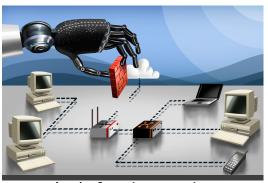


Use Case: Intrusion Response

- A defender owns an infrastructure
 - Consists of connected components
 - Components run network services
 - Defender defends the infrastructure by monitoring and active defense
 - Has partial observability
- An attacker seeks to intrude on the infrastructure
 - Has a partial view of the infrastructure
 - Wants to compromise specific components
 - Attacks by reconnaissance, exploitation and pivoting



Automated Intrusion Response



Levels of security automation







Audit logs
Manual detection.
Manual prevention.



Partial automation.

Manual configuration.

Intrusion detection systems.

Intrusion prevention systems.

High automation.System automatically updates itself.

Research

1980s

1990s

2000s-Now

2/4

Automated Intrusion Response



Can we find effective security strategies through decision-theoretic methods?

Levels of security automation



No automation.

Manual detection

Manual prevention

Operator assistance.
Audit logs
Manual detection.

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Partial automation.

Manual configuration.

Intrusion detection systems

ntrusion prevention systems

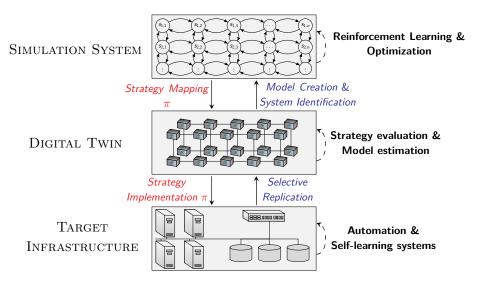
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updates itse

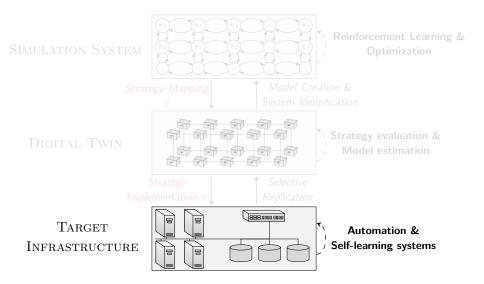
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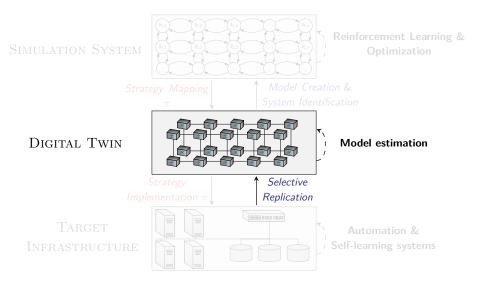
2000s-Now

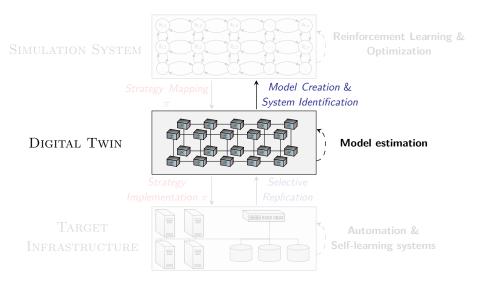
Research

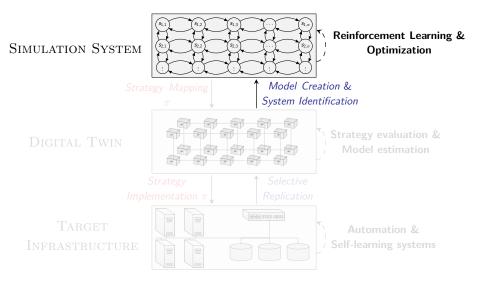


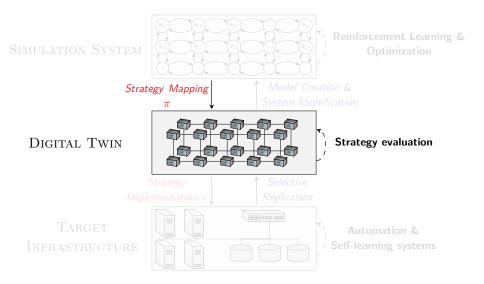


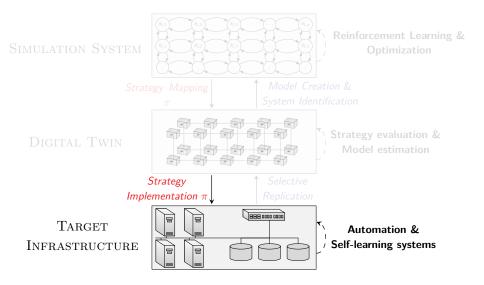


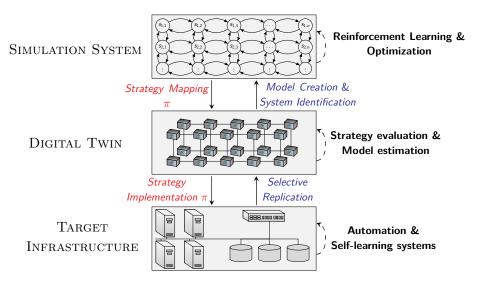




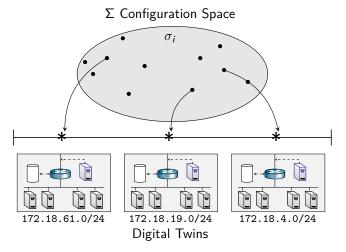








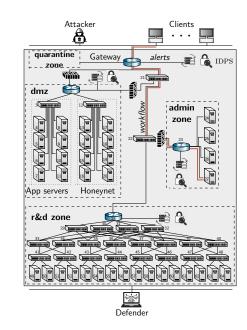
Creating a Digital Twin of the Target Infrastructure



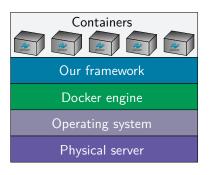
- Given an infrastructure configuration, our framework automates the creation of a digital twin.
- ► The configuration space defines the class of infrastructures that we can emulate.

Example Infrastructure Configuration

- ► 64 nodes
 - 24 OVS switches
 - 3 gateways
 - 6 honeypots
 - 8 application servers
 - 4 administration servers
 - 15 compute servers
- ▶ 11 vulnerabilities
 - ► CVE-2010-0426
 - ► CVE-2015-3306
 - etc.
- Management
 - ▶ 1 SDN controller
 - 1 Kafka server
 - 1 elastic server

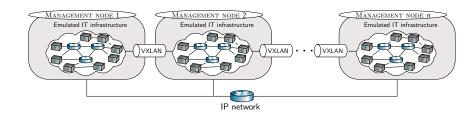


Emulating Physical Components



- We emulate physical components with Docker containers
- Focus on linux-based systems
- Our framework provides the orchestration layer

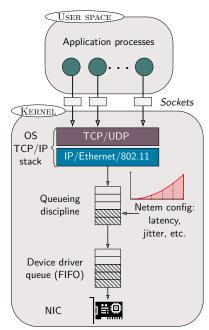
Emulating Network Connectivity



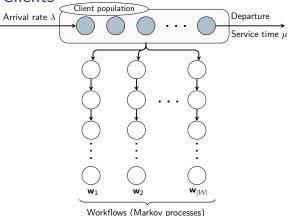
- We emulate network connectivity on the same host using network namespaces
- Connectivity across physical hosts is achieved using VXLAN tunnels with Docker swarm

Emulating Network Conditions

- ► Traffic shaping using NetEm
- ► Allows to configure:
 - Delay
 - Capacity
 - Packet Loss
 - Jitter
 - Queueing delays
 - etc.



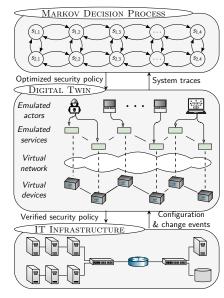
Emulating Clients



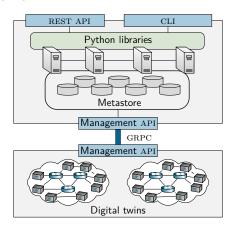
- ► Homogeneous client population
- ▶ Clients arrive according to $Po(\lambda)$
- ightharpoonup Client service times $Exp(\mu)$
- ▶ Service dependencies $(S_t)_{t=1,2,...}$ ~ MC

Emulating The Attacker and The Defender

- API for automated defender and attacker actions
- Attacker actions:
 - Exploits
 - Reconnaissance
 - Pivoting
 - etc.
- Defender actions:
 - Shut downs
 - Redirect
 - Isolate
 - Recover
 - Migrate
 - etc.



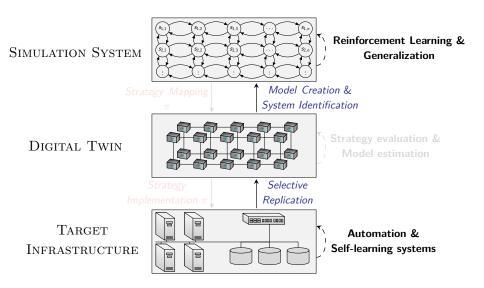
Software framework



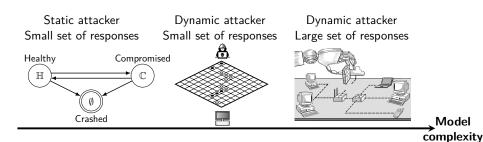
- More details about the software framework
 - ► Source code: https://github.com/Limmen/csle
 - ► Documentation: http://limmen.dev/csle/
 - ▶ Demo: https://www.youtube.com/watch?v=iE2KPmtIs2A
 - Installation:

https://www.youtube.com/watch?v=l_g3sRJwwhc

System Identification

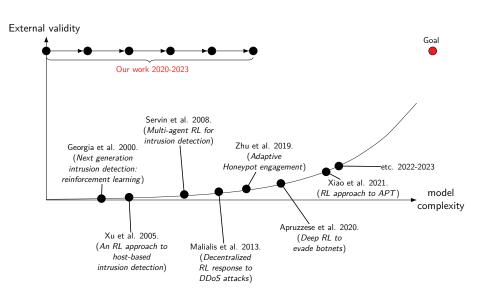


System Model



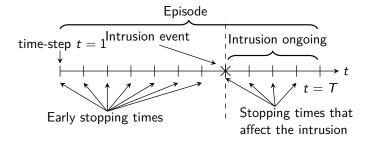
- Intrusion response can be modeled in many ways
 - As a parametric optimization problem
 - As an optimal stopping problem
 - As a dynamic program
 - As a game
 - etc.

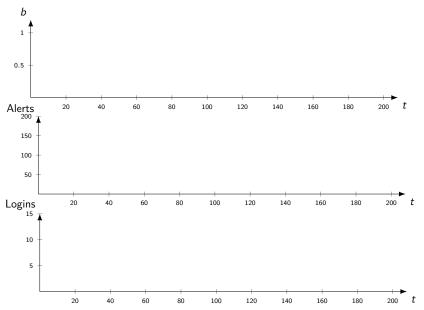
Related Work on Learning Automated Intrusion Response

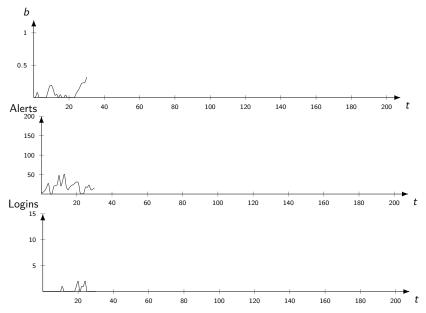


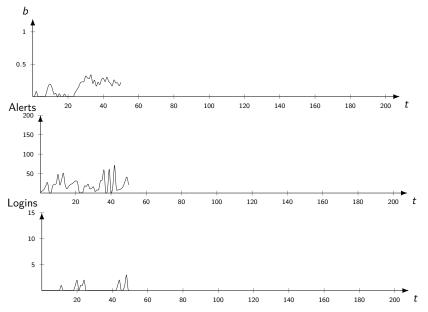
Intrusion Response through Optimal Stopping

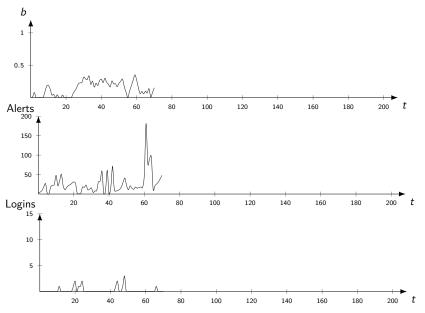
- Suppose
 - The attacker follows a fixed strategy (no adaptation)
 - ▶ We only have one response action, e.g., block the gateway
- Formulate intrusion response as optimal stopping

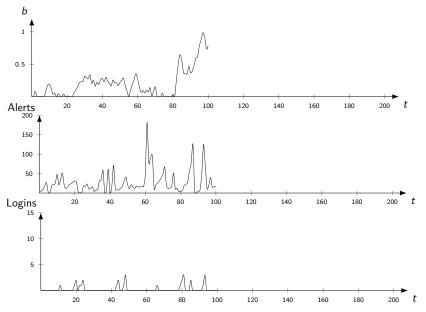


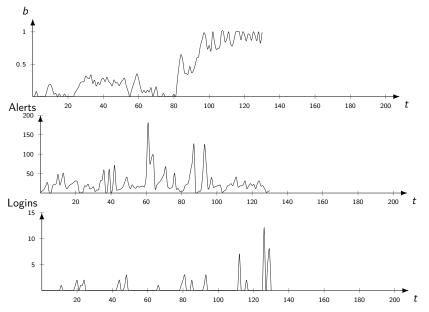


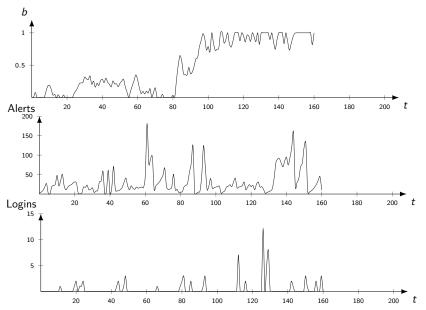


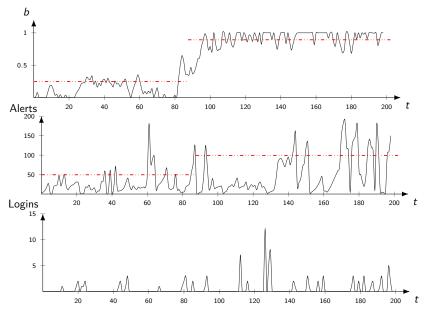


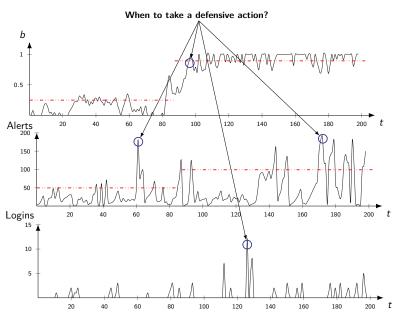












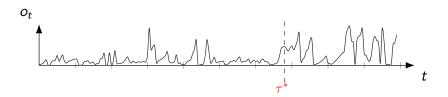
The Defender's Optimal Stopping Problem (1/3)

- ▶ Infrastructure is a **discrete-time dynamical system** $(s_t)_{t=1}^T$
- ▶ Defender observes a **noisy observation process** $(o_t)_{t=1}^T$
- ▶ Two options at each time t: (\mathfrak{C})ontinue and (\mathfrak{S})stop
- ▶ Find the *optimal stopping time* τ^* :

$$\tau^{\star} \in \operatorname*{arg\,max}_{\tau} \mathbb{E}_{\tau} \bigg[\sum_{t=1}^{\tau-1} \gamma^{t-1} \mathcal{R}^{\mathfrak{C}}_{s_{t}s_{t+1}} + \gamma^{\tau-1} \mathcal{R}^{\mathfrak{S}}_{s_{\tau}s_{\tau}} \bigg]$$

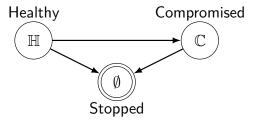
where $\mathcal{R}_{ss'}^{\mathfrak{S}}$ & $\mathcal{R}_{ss'}^{\mathfrak{C}}$ are the stop/continue rewards and au is

$$\tau = \inf\{t : t > 0, a_t = \mathfrak{S}\}\$$



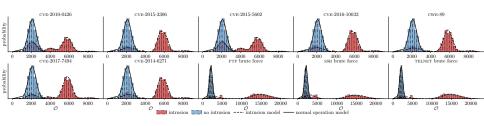
The Defender's Optimal Stopping Problem (2/3)

- ▶ **Objective:** stop the attack as soon as possible
- ▶ Let the **state space** be $S = {\mathbb{H}, \mathbb{C}, \emptyset}$



The Defender's Optimal Stopping Problem (3/3)

Let the observation process $(o_t)_{t=1}^T$ represent IDS alerts



- ► Estimate the observation distribution based on *M* samples from the twin
- ▶ E.g., compute empirical distribution \hat{Z} as estimate of Z
- $ightharpoonup \widehat{Z}
 ightarrow^{\mathsf{a.s}} Z$ as $M
 ightarrow \infty$ (Glivenko-Cantelli theorem)

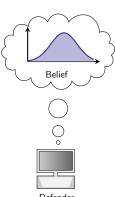
Optimal Stopping Strategy

► The defender can compute the **belief**

$$b_t \triangleq \mathbb{P}[S_t = \mathbb{C} \mid b_1, o_1, o_2, \dots o_t]$$

Stopping strategy:

$$\pi(b): [0,1] \to \{\mathfrak{S},\mathfrak{C}\}$$





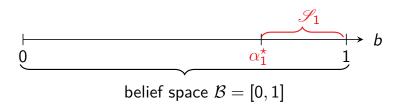
Optimal Threshold Strategy

Theorem

There exists an optimal defender strategy of the form:

$$\pi^*(b) = \mathfrak{S} \iff b \ge \alpha^* \qquad \qquad \alpha^* \in [0, 1]$$

i.e., the stopping set is $\mathscr{S} = [\alpha^*, 1]$

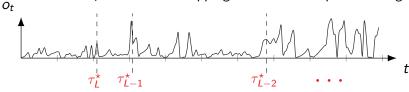


Optimal Multiple Stopping

- ▶ Suppose the defender can take $L \ge 1$ response actions
- ▶ Find the *optimal stopping times* $\tau_L^{\star}, \tau_{L-1}^{\star}, \dots, \tau_1^{\star}$:

$$\begin{split} &(\tau_{l}^{\star})_{l=1,\dots,L} \in \operatorname*{arg\,max}_{\tau_{1},\dots,\tau_{L}} \mathbb{E}_{\tau_{1},\dots,\tau_{L}} \bigg[\sum_{t=1}^{\tau_{L}-1} \gamma^{t-1} \mathcal{R}^{\mathfrak{C}}_{s_{t}s_{t+1}} + \gamma^{\tau_{L}-1} \mathcal{R}^{\mathfrak{S}}_{s_{\tau_{L}}s_{\tau_{L}}} + \\ &\sum_{t=\tau_{L}+1}^{\tau_{L-1}-1} \gamma^{t-1} \mathcal{R}^{\mathfrak{C}}_{s_{t}s_{t+1}} + \gamma^{\tau_{L-1}-1} \mathcal{R}^{\mathfrak{S}}_{s_{\tau_{L-1}}s_{\tau_{L-2}}} + \dots + \\ &\sum_{t=\tau_{2}+1}^{\tau_{1}-1} \gamma^{t-1} \mathcal{R}^{\mathfrak{C}}_{s_{t}s_{t+1}} + \gamma^{\tau_{1}-1} \mathcal{R}^{\mathfrak{S}}_{s_{\tau_{1}}s_{\tau_{1}}} \bigg] \end{split}$$

where τ_I denotes the stopping time with I stops remaining.



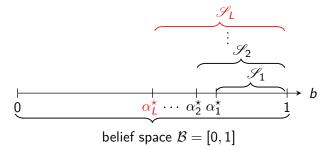
Optimal Multi-Threshold Strategy

Theorem

- ▶ Stopping sets are nested $\mathcal{S}_{l-1} \subseteq \mathcal{S}_l$ for l = 2, ... L.
- ▶ If $(o_t)_{t\geq 1}$ is totally positive of order 2 (TP2), there exists an optimal defender strategy of the form:

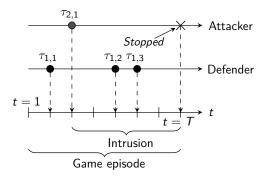
$$\pi_I^*(b) = \mathfrak{S} \iff b \ge \alpha_I^*, \qquad I = 1, \dots, L$$

where $\alpha_I^{\star} \in [0,1]$ is decreasing in I.



Optimal Stopping Game

Suppose the attacker is dynamic and decides when to start and abort its intrusion.



► Find the *optimal stopping times*

$$\max_{\tau_{\mathrm{D},1},\ldots,\tau_{\mathrm{D},L}} \min_{\tau_{\mathrm{A},1},\tau_{\mathrm{A},2}} \mathbb{E}[J]$$

where J is the defender's objective.

Best-Response Multi-Threshold Strategies (1/2)

Theorem

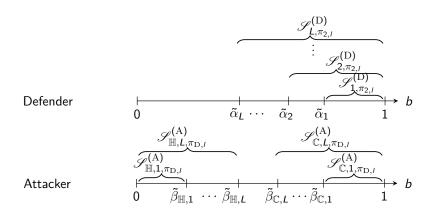
► The defender's best response is of the form:

$$\tilde{\pi}_{\mathrm{D},l}(b) = \mathfrak{S} \iff b \geq \tilde{\alpha}_{l}, \qquad l = 1, \ldots, L$$

► The attacker's best response is of the form:

$$ilde{\pi}_{\mathrm{A},I}(b) = \mathfrak{C} \iff ilde{\pi}_{\mathrm{D},I}(\mathfrak{S} \mid b) \geq ilde{eta}_{\mathbb{H},I}, \quad I = 1,\ldots,L, s = \mathbb{H} \\ ilde{\pi}_{\mathrm{A},I}(b) = \mathfrak{S} \iff ilde{\pi}_{\mathrm{D},I}(\mathfrak{S} \mid b) \geq ilde{eta}_{\mathbb{C},I}, \quad I = 1,\ldots,L, s = \mathbb{C}$$

Best-Response Multi-Threshold Strategies (2/2)



Efficient Computation of Best Responses

Algorithm 1: Threshold Optimization

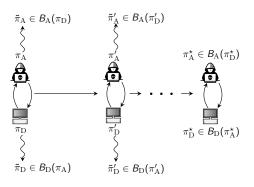
- 1 **Input:** Objective function J, number of thresholds L, parametric optimizer PO
- 2 **Output:** A approximate best response strategy $\hat{\pi}_{\theta}$
- 3 Algorithm

```
4 \Theta \leftarrow [0,1]^L

5 For each \theta \in \Theta, define \pi_{\theta}(b_t) as
6 \pi_{\theta}(b_t) \triangleq \begin{cases} \mathfrak{S} & \text{if } b_t \geq \theta_i \\ \mathfrak{C} & \text{otherwise} \end{cases}
7 J_{\theta} \leftarrow \mathbb{E}_{\pi_{\theta}}[J]
8 \hat{\pi}_{\theta} \leftarrow \mathrm{PO}(\Theta, J_{\theta})
9 return \hat{\pi}_{\theta}
```

Examples of parameteric optimization algorithms: CEM, BO, CMA-ES, DE, SPSA, etc.

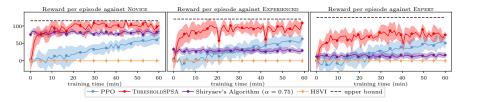
Threshold-Fictitious Play to Approximate an Equilibrium

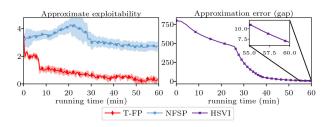


Fictitious play: iterative averaging of best responses.

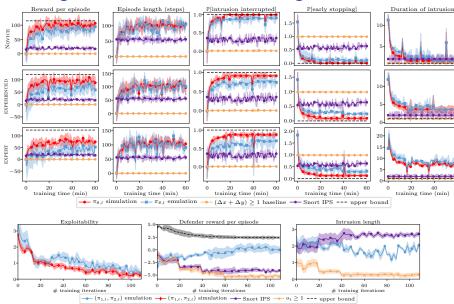
- ▶ Learn best response strategies iteratively
- Average best responses to approximate the equilibrium

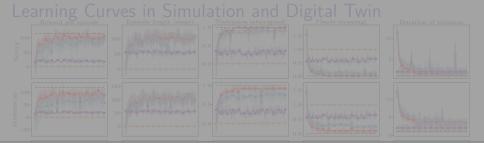
Comparison against State-of-the-art Algorithms



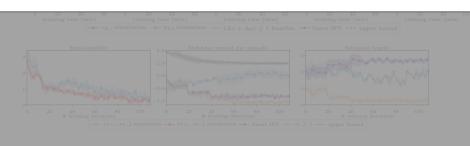


Learning Curves in Simulation and Digital Twin



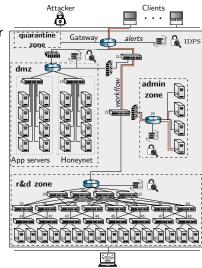


Stopping is about **timing**; now we consider **timing** + action selection



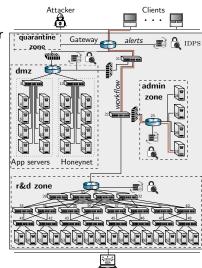
General Intrusion Response Game

- ► Suppose the defender and the attacker can take *L* actions **per node**
- ▶ $\mathcal{G} = \langle \{gw\} \cup \mathcal{V}, \mathcal{E} \rangle$: directed tree representing the virtual infrastructure
- \triangleright \mathcal{V} : set of virtual nodes
- \triangleright \mathcal{E} : set of node dependencies
- \triangleright \mathcal{Z} : set of zones



General Intrusion Response Game

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State Space

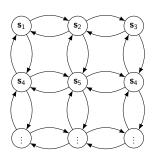
ightharpoonup Each $i \in \mathcal{V}$ has a state

$$\mathbf{v}_{i,t} = (\underbrace{v_{t,i}^{(\mathrm{Z})}}_{\mathrm{D}}, \underbrace{v_{t,i}^{(\mathrm{I})}, v_{t,i}^{(\mathrm{R})}}_{\mathrm{A}})$$

- $lackbox{ System state } \mathbf{s}_t = (\mathbf{v}_{t,i})_{i \in \mathcal{V}} \sim \mathbf{S}_t$
- Markovian time-homogeneous dynamics:

$$\mathbf{s}_{t+1} \sim f(\cdot \mid \mathbf{S}_t, \mathbf{A}_t)$$

$$\mathbf{A}_t = (\mathbf{A}_t^{(\mathrm{A})}, \mathbf{A}_t^{(\mathrm{D})})$$
 are the actions.



State Space

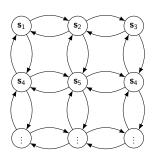
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State Space

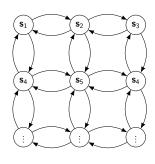
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$$\mathbf{s}_{t+1} \sim f(\cdot \mid \mathbf{S}_t, \mathbf{A}_t)$$

 $\mathbf{A}_t = (\mathbf{A}_t^{(\mathrm{A})}, \mathbf{A}_t^{(\mathrm{D})})$ are the actions.



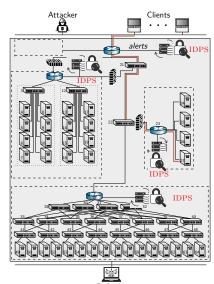
Observations

► IDPSs inspect network traffic and generate alert vectors:

$$\mathbf{o}_t \triangleq \left(\mathbf{o}_{t,1}, \dots, \mathbf{o}_{t,|\mathcal{V}|}
ight) \in \mathbb{N}_0^{|\mathcal{V}|}$$

 $\mathbf{o}_{t,i}$ is the number of alerts related to node $i \in \mathcal{V}$ at time-step t.

 $\mathbf{o}_t = (\mathbf{o}_{t,1}, \dots, \mathbf{o}_{t,|\mathcal{V}|})$ is a realization of the random vector \mathbf{O}_t with joint distribution Z



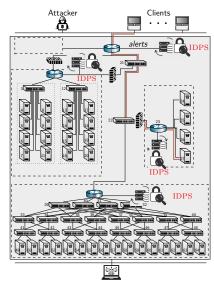
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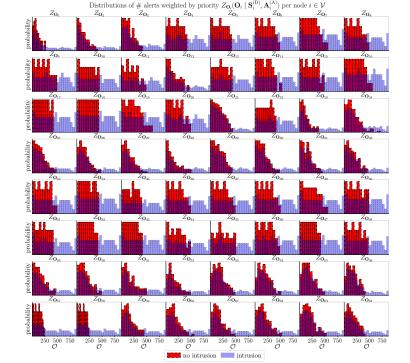
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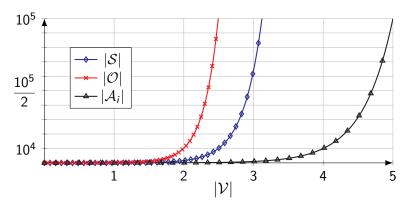
The (General) Intrusion Response Problem

$$\mathsf{maximize}_{\pi_{\mathrm{D}} \in \Pi_{\mathrm{D}}} \ \mathsf{\underline{minimize}}_{\pi_{\mathrm{A}} \in \Pi_{\mathrm{A}}} \ \mathbb{E}_{(\pi_{\mathrm{D}}, \pi_{\mathrm{A}})} \left[J \right]$$

 $\mathbb{E}_{(\pi_D,\pi_A)}$ denotes the expectation of the random vectors $(\mathbf{S}_t,\mathbf{O}_t,\mathbf{A}_t)_{t\in\{1,\ldots,\mathcal{T}\}}$ when following the strategy profile (π_D,π_A)

The Curse of Dimensionality

Solving the game is computationally intractable. The state, action, and observation spaces of the game **grow** exponentially with $|\mathcal{V}|$.



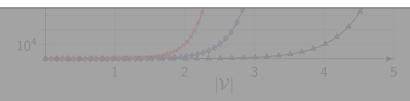
Growth of $|\mathcal{S}|$, $|\mathcal{O}|$, and $|\mathcal{A}_i|$ in function of the number of nodes $|\mathcal{V}|$

The Curse of Dimensionality

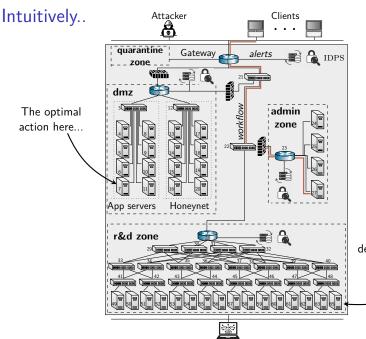
While (1) has a solution (i.e the game Γ has a value (Thm 1)), computing it is intractable since the state, action, and observation spaces of the game **grow exponentially** with $|\mathcal{V}|$



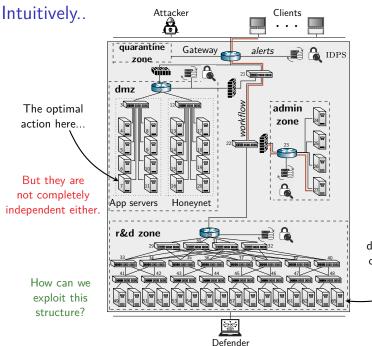
We tackle the scability challenge with **decomposition**



Growth of $|\mathcal{S}|$, $|\mathcal{O}|$, and $|\mathcal{A}_i|$ in function of the number of nodes $|\mathcal{V}|$

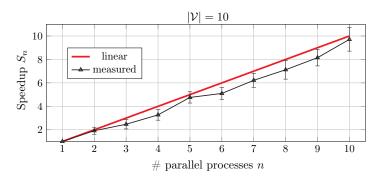


Does not directly depend on the state or action of a node down here



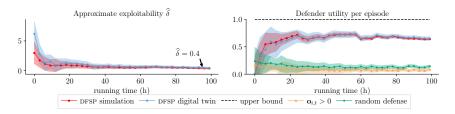
Does not directly depend on the state or action of a node down here

Scalable Learning through Decomposition



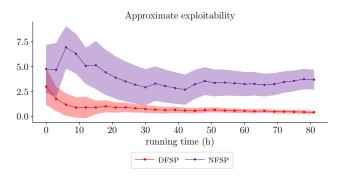
Speedup of best response computation for the decomposed game; T_n denotes the completion time with n processes; the speedup is calculated as $S_n = \frac{T_1}{T_n}$; the error bars indicate standard deviations from 3 measurements.

Learning Equilibrium Strategies



Learning curves obtained during training of DFSP to find optimal (equilibrium) strategies in the intrusion response game; **red and blue curves relate to dfsp**; black, orange and green curves relate to baselines.

Comparison with NFSP



Learning curves obtained during training of DFSP and NFSP to find optimal (equilibrium) strategies in the intrusion response game; **the red curve relate to dfsp** and the purple curve relate to NFSP; all curves show simulation results.

Conclusions

- We develop a framework to automatically learn security strategies.
- We apply the framework to an intrusion response use case.
- We derive properties of optimal security strategies.
- We evaluate strategies on a digital twin.
- ▶ Questions → demonstration

