Intrusion Prevention through Optimal Stopping and Self-Play NSE Seminar

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Use Case: Intrusion Prevention

- A Defender owns an infrastructure
 - Consists of connected components
 - Components run network services
 - Defender defends the infrastructure by monitoring and active defense
 - Has partial observability
- An Attacker seeks to intrude on the infrastructure
 - Has a partial view of the infrastructure
 - Wants to compromise specific components
 - Attacks by reconnaissance, exploitation and pivoting





















3/28





— Reference points

----- Intrusion prevention milestones











Our Approach

Formal model:

- Controlled Hidden Markov Model
- Defender has partial observability
- A game if attacker is active

Data collection:

- Emulated infrastructure
- Finding defender strategies:
 - Self-play reinforcement learning
 - Optimal stopping



Use Case & Approach:

- Intrusion prevention
- Reinforcement learning and optimal stopping
- Formal Model of The Use Case
 - Intrusion prevention as an optimal stopping problem
 - Partially observed stochastic game

▶ Game Analysis and Structure of (π
₁, π
₂)
 ▶ Structural result: multi-threshold best responses
 ▶ Stopping sets S_l⁽¹⁾ are connected and nested

Our Method

- Emulation of the target infrastructure,
- System identification
- Our reinforcement learning algorithm
- Results & Conclusion
 - Numerical evaluation results & Demo
 - Conclusion & future work

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- Numerical evaluation results & Demo
- Conclusion & future work

- The system evolves in discrete time-steps.
- ▶ A stop action = a defensive action. (e.g. reconfigure IPS)



• The L - Ith stopping time τ_I is:

 $\tau_{l} = \inf\{t : t > \tau_{l-1}, a_{t} = S\}, \qquad l \in 1, .., L, \ \tau_{L+1} = 0$

► τ_l is a random variable from sample space Ω to \mathbb{N} , which is dependent on $h_{\tau} = \rho_1, a_1, o_1, \ldots, a_{\tau-1}, o_{\tau}$ and independent of $a_{\tau}, o_{\tau+1}, \ldots$

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We consider the class of stopping times $\mathcal{T}_t = \{\tau_l \leq t | \tau_l > \tau_{l-1}\} \in \mathcal{F}_k \ (\mathcal{F}_k = natural filtration on h_t).$







The Defender's Stop Actions



Infrastructure

- 1. Ingress traffic goes through deep packet inspection at gateway
- 2. Gateway runs the Snort IDS/IPS and may drop packets
- 3. The defender controls the IPS configuration
- 4. At each stopping time, we update the IPS configuration

The Defender's Stop Actions

T	Alarm class c	Action
34	Attempted administrator privilege gain	DROP
33	Attempted user privilege gain	DROP
32	Inappropriate content was detected	DROP
31	Potential corporate privacy violation	DROP
30	Executable code was detected	DROP
29	Successful administrator privilege gain	DROP
28	Successful user privilege gain	DROP
27	A network trojan was detected	DROP
26	Unsuccessful user privilege gain	DROP
25	Web application attack	DROP
24	Attempted denial of service	DROP
23	Attempted information leak	DROP
22	Potentially Bad Traffic	DROP
21	Attempt to login by a default username and password	DROP
20	Detection of a denial of service attack	DROP
:	:	:

Table 1: Defender stop actions in the emulation; *I* denotes the number of stops remaining.

- We assume a strategic attacker
 - Attacker knows which actions generate alarms
 - Attacker tries to be stealthy
 - Attacker may try to achieve denial of service





The attacker's stopping times τ_{2,1}, τ_{2,2},... determine the times to start/stop the intrusion

- During the inrusion, the attacker follows a fixed intrusion strategy
- ▶ The defender's stopping times $\tau_{1,L}, \tau_{1,L-1}, \ldots$ determine the times to update the IPS configuration



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- ► The defender's stopping times \u03c6_{1,L}, \u03c6₁



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 - During the intrusion, the attacker follows a fixed intrusion strategy
- The defender's stopping times τ_{1,1}, τ_{1,2},... determine the times to update the IPS configuration
- We seek a Nash equilibrium (π₁^{*}, π₂^{*}), from which we can extract the optimal defender strategy π₁^{*} against a worst-case attacker.
Optimal Stopping Game



We model this game as a zero-sum partially observed stochastic game

- The attacker's stopping times \(\tau_{2,1}, \tau_{2,2}, \ldots\) determine the times to start/stop the intrusion
 - During the intrusion, the attacker follows a fixed intrusion strategy
- The defender's stopping times τ_{1,1}, τ_{1,2},... determine the times to update the IPS configuration
- We seek a **Nash equilibrium** (π_1^*, π_2^*) , from which we can extract the optimal defender strategy π_1^* against the

• Players: $\mathcal{N} = \{1, 2\}$ (Defender=1)

States: Intrusion $s_t \in \{0, 1\}$, terminal \emptyset .

Observations:

- ► IDS Alerts $\Delta x_{1,t}, \Delta x_{2,t}, \dots, \Delta x_{M,t}$, defender stops remaining $l_t \in \{1, ..., L\}$, $f_X(\Delta x_1, \dots, \Delta x_M | s_t)$
- Actions:

•
$$\mathcal{A}_1 = \mathcal{A}_2 = \{S, C\}$$

Rewards:

- Defender reward: security and service.
- Attacker reward: negative of defender reward.
- Transition probabilities:
 - Follows from game dynamics.



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One-Sided Partial Observability

- We assume that the attacker has perfect information. Only the defender has partial information.
- The attacker's view:



The defender's view:



One-Sided Partial Observability

- We assume that the attacker has perfect information. Only the defender has partial information.
- The attacker's view:



► The defender's view:



• Makes it tractable to compute the defender's belief $b_{t,\pi_2}^{(1)}(s_t) = \mathbb{P}[s_t|h_t,\pi_2]$ (avoid nested beliefs)

- Defender strategy is of the form: $\pi_{1,l} : \mathcal{B} \to \Delta(\mathcal{A}_1)$
- Attacker strategy is of the form: $\pi_{2,l} : S \times B \to \Delta(A_2)$

Defender and attacker objectives:

$$J_{1}(\pi_{1,l},\pi_{2,l}) = \mathbb{E}_{(\pi_{1,l},\pi_{2,l})} \left[\sum_{t=1}^{\infty} \gamma^{t-1} \mathcal{R}_{l}(s_{t}, \boldsymbol{a}_{t}) \right]$$
$$J_{2}(\pi_{1,l},\pi_{2,l}) = -J_{1}$$

Best response correspondences:

$$B_1(\pi_{2,l}) = \arg\max_{\substack{\pi_{1,l} \in \Pi_1 \\ \pi_{2,l} \in \Pi_2}} J_1(\pi_{1,l}, \pi_{2,l})$$
$$B_2(\pi_{1,l}) = \arg\max_{\substack{\pi_{2,l} \in \Pi_2}} J_2(\pi_{1,l}, \pi_{2,l})$$

Nash equilibrium $(\pi_{1,l}^*, \pi_{2,l}^*)$:

$$\pi^*_{1, \textit{l}} \in B_1(\pi^*_{2, \textit{l}})$$
 and $\pi^*_{2, \textit{l}} \in B_2(\pi^*_{1, \textit{l}})$

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Nash equilibrium $(\pi_{1,l}^*, \pi_{2,l}^*)$:

 $\pi_{1,l}^* \in B_1(\pi_{2,l}^*)$ and $\pi_{2,l}^* \in B_2(\pi_{1,l}^*)$

- Defender strategy is of the form: $\pi_{1,I} : \mathcal{B} \to \Delta(\mathcal{A}_1)$
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Defender and attacker objectives:

$$J_{1}(\pi_{1,l},\pi_{2,l}) = \mathbb{E}_{(\pi_{1,l},\pi_{2,l})} \left[\sum_{t=1}^{\infty} \gamma^{t-1} \mathcal{R}_{l}(s_{t}, \boldsymbol{a}_{t}) \right]$$
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▶ Nash equilibrium $(\pi_{1,l}^*, \pi_{2,l}^*)$: $\pi_{1,l}^* \in B_1(\pi_{2,l}^*)$ and $\pi_{2,l}^* \in B_2(\pi_{1,l}^*)$











Attacker



















Emulating the Target Infrastructure

- Emulate hosts with docker containers
- Emulate IDS and vulnerabilities with software
- Network isolation and traffic shaping through NetEm in the Linux kernel
- Enforce resource constraints using cgroups.
- Emulate client arrivals with Poisson process
- Internal connections are full-duplex & loss-less with bit capacities of 1000 Mbit/s
- External connections are full-duplex with bit capacities of 100 Mbit/s & 0.1% packet loss in normal operation and random bursts of 1% packet loss



Running a Game Episode in the Emulation

- A distributed system with synchronized clocks
- We run software sensors on all emulated hosts
- Sensors produce messages to a distributed queue (Kafka)
- A stream processor (Spark) consumes messages from the queue and computes statistics
- Actions are selected based on the computed statistics and the strategies
- Actions are sent to the emulation using gRPC
- Actions are executed by running commands on the hosts



Our Reinforcement Learning Approach

- ► We learn a Nash equilibrium $(\pi_{1,l,\theta^{(1)}}^*, \pi_{2,l,\theta^{(2)}}^*)$ through fictitious self-play.
- In each iteration:
 - 1. Learn a best response strategy of the defender by solving a POMDP $\tilde{\pi}_{1,I,\theta^{(1)}} \in B_1(\pi_{2,I,\theta^{(2)}})$.
 - 2. Learn a best response strategy of the attacker by solving an MDP $\tilde{\pi}_{2,I,\theta^{(2)}} \in B_2(\pi_{1,I,\theta^{(1)}}).$
 - 3. Store the best response strategies in two buffers Θ_1, Θ_2
 - 4. Update strategies to be the average of the stored strategies



Input

 Γ^{p} : the POSG

Output

 $(\pi^*_{1, heta},\pi^*_{2, heta})$: an approximate Nash equilibrium

1: procedure APPROXIMATEFP

$$\begin{array}{l} 2: \qquad \theta^{(1)} \sim \mathcal{N}_{L}(-1,1), \quad \theta^{(2)} \sim \mathcal{N}_{2L}(-1,1) \\ \Omega^{(1)} \sim \mathcal{N}_{L}(0,1), \quad \Omega^{(2)} \sim \mathcal{N}_{2L}(-1,1) \end{array}$$

$$\Theta^{(1)} \leftarrow \{\theta^{(1)}\}, \quad \Theta^{(2)} \leftarrow \{\theta^{(2)}\}, \quad \delta \leftarrow \infty$$

4: while
$$\hat{\delta} \ge \delta$$
 do

5:
$$\theta^{(1)} \leftarrow \text{THRESHOLDBR}(\Gamma^{p}, \pi_{2,l,\theta}, N, a, c, \lambda, A, \epsilon)$$

6:
$$\theta^{(2)} \leftarrow \text{THRESHOLDBR}(\Gamma^{p}, \pi_{1,l,\theta}, N, a, c, \lambda, A, \epsilon)$$

$$7: \qquad \Theta^{(1)} \leftarrow \Theta^{(2)} \cup \theta^{(1)}, \quad \Theta_2 \leftarrow \Theta^{(2)} \cup \theta^{(2)}$$

- 8: $\pi_{1,l,\theta} \leftarrow \text{MIXTUREDISTRIBUTION}(\Theta^{(1)})$
- 9: $\pi_{2,l,\theta} \leftarrow \text{MIXTUREDISTRIBUTION}(\Theta^{(2)})$

10:
$$\hat{\delta} = \text{EXPLOITABILITY}(\pi_{1,l,\theta}, \pi_{2,l,\theta})$$

11: end while

12: **return**
$$(\pi_{1,l,\theta}, \pi_{2,l,\theta})$$

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- while $\hat{\delta} > \delta$ do 4.
- $\theta^{(1)} \leftarrow \text{THRESHOLDBR}(\Gamma^{p}, \pi_{2,l,\theta}, N, a, c, \lambda, A, \epsilon)$ 5.
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- $\pi_{1,l,\theta} \leftarrow \text{MIXTUREDISTRIBUTION}(\Theta^{(1)})$ 8:
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- $\hat{\delta} = \text{EXPLOITABILITY}(\pi_{1,l,\theta}, \pi_{2,l,\theta})$ 10.
- end while 11:
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- 13: end procedure

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10:
$$\hat{\delta} = ext{ExploitAbility}(\pi_{1,l, heta},\pi_{2,l, heta})$$

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1: procedure APPROXIMATEFP

$$\begin{array}{ccc} & & \theta^{(1)} \sim \mathcal{N}_{L}(-1,1), & \theta^{(2)} \sim \mathcal{N}_{2L}(-1,1) \\ & & & & \\ & & & & \\ & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & & \\ & & & \\ & & & & \\ & & & & \\ & & & & \\ & & &$$

$$\Theta^{(1)} \leftarrow \{\theta^{(1)}\}, \quad \Theta^{(2)} \leftarrow \{\theta^{(2)}\}, \quad \hat{\delta} \leftarrow \infty$$

4: while
$$\hat{\delta} \ge \delta$$
 do

5:
$$\theta^{(1)} \leftarrow \text{THRESHOLDBR}(\Gamma^{p}, \pi_{2,l,\theta}, N, a, c, \lambda, A, \epsilon)$$

6:
$$\theta^{(2)} \leftarrow \text{THRESHOLDBR}(\Gamma^{p}, \pi_{1,l,\theta}, N, a, c, \lambda, A, \epsilon)$$

$$7: \qquad \Theta^{(1)} \leftarrow \Theta^{(2)} \cup \theta^{(1)}, \quad \Theta_2 \leftarrow \Theta^{(2)} \cup \theta^{(2)}$$

- 8: $\pi_{1,l,\theta} \leftarrow \text{MIXTUREDISTRIBUTION}(\Theta^{(1)})$
- 9: $\pi_{2,l,\theta} \leftarrow \text{MIXTUREDISTRIBUTION}(\Theta^{(2)})$

$$\hat{\delta} = ext{Exploitability}(\pi_{1,l, heta},\pi_{2,l, heta})$$

11: end while

10:

12: **return**
$$(\pi_{1,l,\theta}, \pi_{2,l,\theta})$$
Our Reinforcement Learning Algorithm for Learning Best-Response Threshold Strategies

- We use the structural result that threshold best response strategies exist (Theorem 1) to design an efficient reinforcement learning algorithm to learn best response strategies.
- We seek to learn:
 - *L* thresholds of the defender, $\tilde{\alpha}_1, \geq \tilde{\alpha}_2, \ldots, \geq \tilde{\alpha}_L \in [0, 1]$
 - ▶ 2*L* thresholds of the attacker, $\tilde{\beta}_{0,1}, \tilde{\beta}_{1,1}, \dots, \tilde{\beta}_{0,L}, \tilde{\beta}_{1,L} \in [0,1]$

We learn these thresholds iteratively through Robbins and Monro's stochastic approximation algorithm.¹



¹Herbert Robbins and Sutton Monro. "A Stochastic Approximation Method". In: The Annals of Mathematical Statistics 22.3 (1951), pp. 400–407. DOI: 10.1214/aoms/1177729586. URL: https://doi.org/10.1214/aoms/1177729586.

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- 1. Parameterize the strategies $\pi_{1,l,\theta^{(1)}}$, $\pi_{1,l,\theta^{(2)}}$ by $\theta^{(1)} \in \mathbb{R}^L$, $\theta^{(2)} \in \mathbb{R}^{2L}$
- 2. The policy gradient

$$\nabla_{\theta^{(i)}} J(\theta^{(i)}) = \mathbb{E}_{\pi_{i,l,\theta^{(i)}}} \left[\sum_{t=1}^{\infty} \nabla_{\theta^{(i)}} \log \pi_{i,l,\theta^{(i)}}(a_t^{(i)}|s_t) \sum_{\tau=t}^{\infty} r_t \right]$$

exists as long as $\pi_{i,l,\theta^{(i)}}$ is differentiable.

3. A pure threshold strategy is not differentiable.

4. To ensure differentiability and to constrain the thresholds to be in [0, 1], we define $\pi_{i,\theta^{(l)},l}$ to be a smooth stochastic strategy that approximates a threshold strategy:

$$\pi_{i,\theta^{(i)}}(S|b(1)) = \left(1 + \left(\frac{b(1)(1 - \sigma(\theta^{(i),j}))}{\sigma(\theta^{(i),j})(1 - b(1))}\right)^{-20}\right)^{-1}$$

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Smooth Threshold



- 1. We learn the thresholds through simulation.
- 2. For each iteration $n \in \{1, 2, ...\}$, we perturb $\theta_n^{(i)}$ to obtain $\theta_n^{(i)} + c_n \Delta_n$ and $\theta_n^{(i)} c_n \Delta_n$.
- 3. Then, we run two MDP or POMDP episodes
- 4. We then use the obtained episode outcomes $\hat{J}_i(\theta_n^{(i)} + c_n\Delta_n)$ and $\hat{J}_i(\theta_n^{(i)} - c_n\Delta_n)$ to estimate $\nabla_{\theta^{(i)}}J_i(\theta^{(i)})$ using the Simultaneous Perturbation Stochastic Approximation (SPSA) gradient estimator⁴:

$$\left(\hat{\nabla}_{\theta_n^{(i)}} J_i(\theta_n^{(i)})\right)_k = \frac{\hat{J}_i(\theta_n^{(i)} + c_n \Delta_n) - \hat{J}_i(\theta_n^{(i)} - c_n \Delta_n)}{2c_n(\Delta_n)_k}$$

$$\theta_{n+1}^{(i)} = \theta_n^{(i)} + a_n \hat{\nabla}_{\theta_n^{(i)}} J_i(\theta_n^{(i)})$$

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Simulation Results (Emulation Results TBD..)



Demo - A System for Interactive Examination of Learned Security Strategies



Architecture of the system for examining learned security strategies.

Conclusions & Future Work

Conclusions:

We develop a *method* to automatically learn security policies

 (1) emulation system; (2) system identification; (3) simulation system; (4) reinforcement learning and (5) domain randomization and generalization.

We apply the method to an intrusion prevention use case

- We formulate intrusion prevention as a multiple stopping problem
 - We present a Partially Observed Stochastic Game of the use case
 - We present a POMDP model of the defender's problem
 - We present a MDP model of the attacker's problem
 - We apply the stopping theory to establish structural results of the best response strategies

Our research plans:

- Run experiments in the emulation system
- Make learned strategy available as plugin to the Snort IDS
- Extend the model
 - Less restrictions on defender

Scaling up the emulation system:

Non-static infrastructures