

Intrusion Tolerance for Networked Systems through Two-Level Feedback Control

IEEE DSN 2024, Brisbane, Australia
International Conference on Dependable Systems and Networks

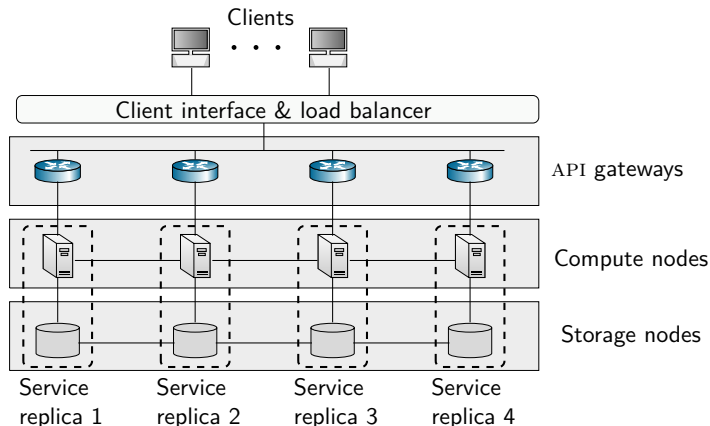
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June 27, 2024

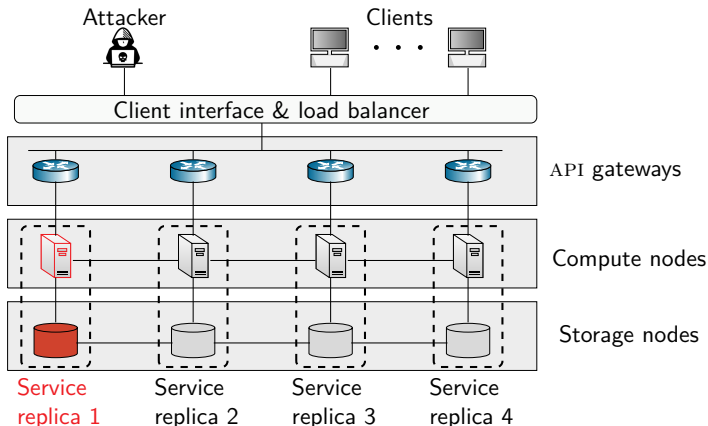


Use Case: Intrusion Tolerance



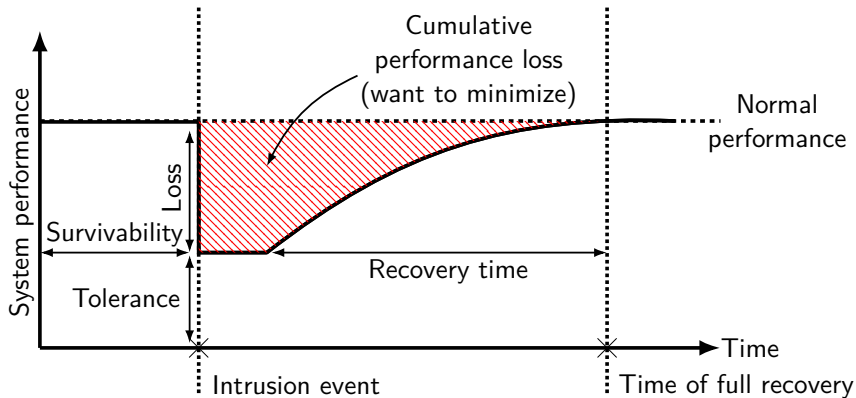
- ▶ A **replicated system** offers a service to a client population.
- ▶ The system should provide **service without disruption**.

Use Case: Intrusion Tolerance



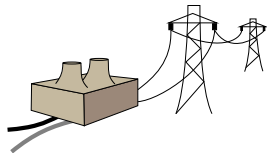
- ▶ An **attacker** seeks to intrude on the system and disrupt service.
- ▶ The system should **tolerate intrusions**.

Intrusion Tolerance (Simplified)



Increasing Demand for Intrusion-Tolerant Systems

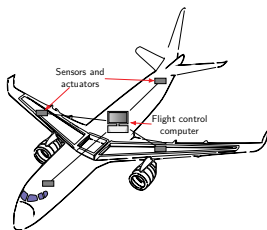
- ▶ As our **reliance on online services grows**, there is an increasing demand for intrusion-tolerant systems.
- ▶ **Example applications:**



Power grids
e.g., SCADA systems¹.



Safety-critical IT systems
e.g., banking systems,
e-commerce applications²,
healthcare systems, etc.



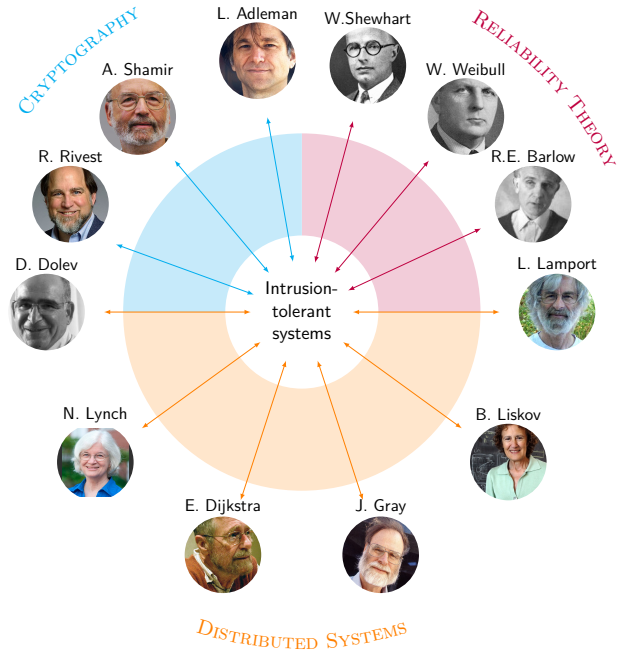
Real-time control systems
e.g., flight control computer³.

¹Amy Babay et al. "Network-Attack-Resilient Intrusion-Tolerant SCADA for the Power Grid". In: *2018 48th Annual IEEE/IFIP International Conference on Dependable Systems and Networks (DSN)*. 2018, pp. 255–266. DOI: [10.1109/DSN.2018.00036](https://doi.org/10.1109/DSN.2018.00036).

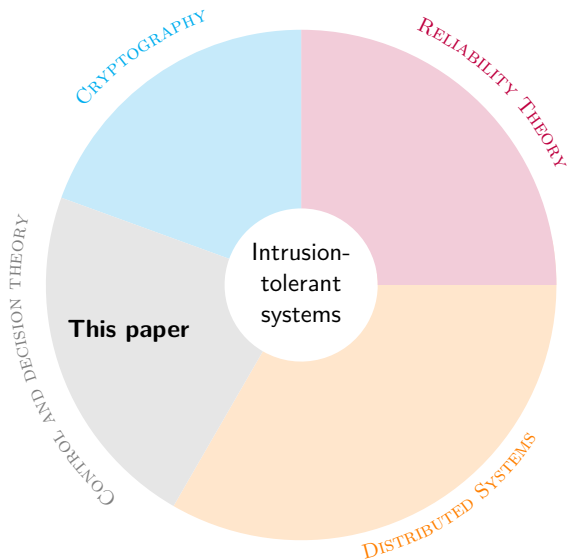
²Jukka Soikkeli et al. "Redundancy Planning for Cost Efficient Resilience to Cyber Attacks". In: *IEEE Transactions on Dependable and Secure Computing* 20.2 (2023), pp. 1154–1168. DOI: [10.1109/TDSC.2022.3151462](https://doi.org/10.1109/TDSC.2022.3151462).

³J.H. Wensley et al. "SIFT: Design and analysis of a fault-tolerant computer for aircraft control". In: *Proceedings of the IEEE* 66.10 (1978), pp. 1240–1255. DOI: [10.1109/PROC.1978.11114](https://doi.org/10.1109/PROC.1978.11114).

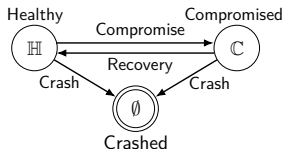
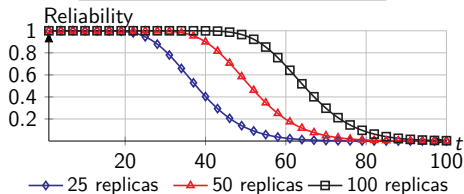
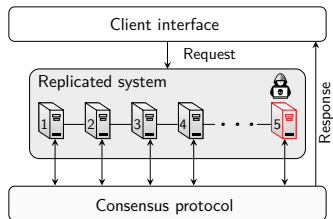
Theoretical Foundations of Intrusion Tolerance



Our Contribution



Building Blocks of An Intrusion-Tolerant System



1. Intrusion-tolerant consensus protocol

A quorum needs to reach agreement to tolerate f compromised replicas.

2. Replication strategy

Cost-reliability trade-off.

3. Recovery strategy

Compromises will occur as $t \rightarrow \infty$.

The Rampart Toolkit for Building High-Integrity Services

Michael K. Reiter

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Abstract. Rampart is a toolkit of protocols to facilitate the implementation of *high-integrity* services, i.e., distributed services that guarantee availability and correctness despite the malicious behavior of some component servers by an attacker. At the core of Rampart are several protocols that solve several basic problems in distributed computing (including asynchronous group membership, reliable broadcast, consensus, agreement), and atomic multicast. Using these protocols, Rampart supports the development of high-integrity services via *machine replication*, and also extends this technique with a new approach to server output voting. In this paper we give a brief overview of Rampart, focusing primarily on its protocol architecture. We also sketch its performance in our prototype implementation and ongoing work.

Published 1995

- Fixed number of replicas
- No recoveries

Prior Work on Intrusion-Tolerant Systems

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The SecureRing Protocols for Securing Group Communication*

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Abstract

The SecureRing group communication protocols provide reliable ordered message delivery and group membership services despite Byzantine faults such as might be caused by modifications to the programs of a group member following illicit access to, or capture of, a group member.

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The approach adopted by SecureRing to protect Byzantine faults is to optimize the performance (fault-free) operation and to pay a performance

Practical Byzantine Fault Tolerance and Proactive Recovery

MIGUEL CASTRO

Microsoft Research

and

BARBARA LISKOV

MIT Laboratory for Computer Science

Published 2002

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- Fixed number of replicas
- **Periodic** recoveries

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Our growing reliance on online services accessible on the Internet demands highly available systems that provide correct service without interruptions. Software bugs, operator mistakes, and malicious attacks are a major cause of service interruptions and they can cause arbitrary behavior that is, Byzantine faults. This article describes a new replication algorithm, BFT, that can be used to build highly available systems that tolerate Byzantine faults. BFT can be used in practice to implement real services: it performs well, it is safe in asynchronous environments such as the Internet, it incorporates mechanisms to defend against Byzantine-faulty clients, and it recovers replicas proactively. The recovery mechanism allows the algorithm to tolerate any number of faults

A Qualitative Analysis of the Intrusion-Tolerance Capabilities of the MAFTIA Architecture

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School of Computing Science, University of Newcastle upon Tyne, UK

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Ian.Welch@mcs.vu

Published 2004

Abstract

MAFTIA was a three-year European research project that explored the use of fault-tolerant techniques to build intrusion-tolerant systems. The MAFTIA architecture embodies a number of key design

- Fixed number of replicas
- Periodic recoveries

Prior Work on Intrusion-Tolerant Systems

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An architecture for adaptive intrusion-tolerant applications

Partha Pal^{1,*} and Paul Rubel¹, Michael Atighetchi¹, Franklin Webber¹, William H. Sanders², Mouna Seri², HariGovind Ramasamy³, James Lyons², Tod Courtney³, Adnan Agbaria², Michel Cukier³, Jeanna Gossett⁴, Idit Keidar⁵

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Published 2006

- Adaptive replication based on heuristics
- Periodic recoveries

Prior Work on Intrusion-Tolerant Systems

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⁵ Department of Electrical Engineering, Technion - Israel Institute of Technology, raik@technion.ac.il

Worm-IT – A wormhole-based intrusion-tolerant group communication system

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Received 26 October 2005; in final form 26 October 2005

Published 2006

- Fixed number of replicas
- Periodic recoveries

Prior Work on Intrusion-Tolerant Systems

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Practical Byzantine Fault Tolerance and Proactive Recovery

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Resilient Intrusion Tolerance through Proactive and Reactive Recovery*

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Published 2007

- Fixed number of replicas
- **Supports both periodic and reactive recoveries**
- Does not provide reactive recovery strategies

Prior Work on Intrusion-Tolerant Systems

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State Transfer for Hypervisor-Based Proactive Recovery of Heterogeneous Replicated Services

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Published 2011

- Fixed number of replicas
- Periodic recoveries

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Network-Attack-Resilient Intrusion-Tolerant SCADA for the Power Grid

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Published 2018

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A Qualitative Analysis of the Intrusion-Tolerance Capabilities of the MAFTIA Architecture

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Worm-IT – A wormhole-based intrusion-tolerant group communication system

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Received 28 October 2001; received in revised form 28 March 2006; accepted 30 March 2006

An architecture for adaptive intrusion-tolerant applications

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Resilient Intrusion Tolerance through Proactive a

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State Transfer for Hypervisor-Based Proactive Recovery of Heterogeneous Replicated Services

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Network-Attack-Resilient Intrusion-Tolerant SCADA for the Power Grid

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Skynet: a Cyber-Aware Intrusion Tolerant Overseer

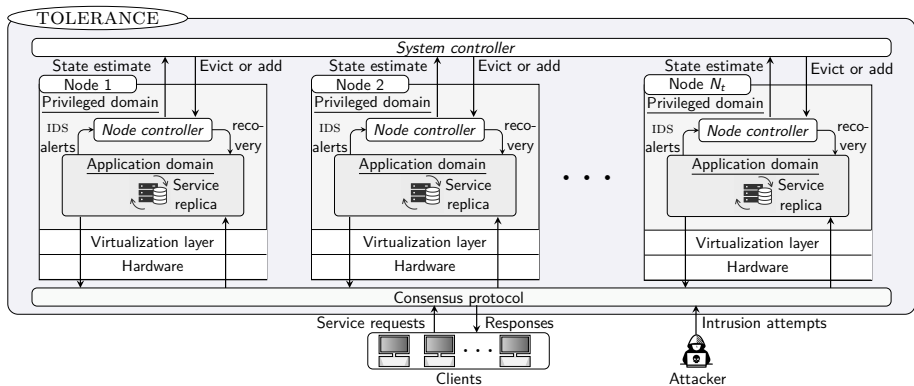
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Published 2023

- Fixed number of replicas
- Periodic recoveries

The TOLERANCE Architecture

Two-level recovery and replication control with feedback.



Definition 1 (Correct service)

The system provides **correct service** if the healthy replicas satisfy the following properties:

- Each request is eventually executed. (Liveness)
- Each executed request was sent by a client. (Validity)
- Each replica executes the same request sequence. (Safety)

Proposition 1 (Correctness of TOLERANCE)

*A system that implements the TOLERANCE architecture **provides correct service** if*

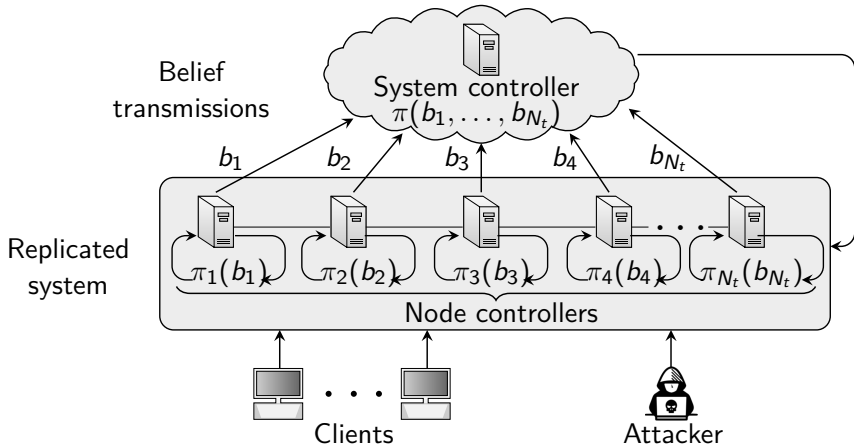
Network links are authenticated.

At most f nodes are compromised or crashed simultaneously.

$N_t \geq 2f + 1$.

The system is partially synchronous.

Intrusion Tolerance as a Two-Level Control Problem



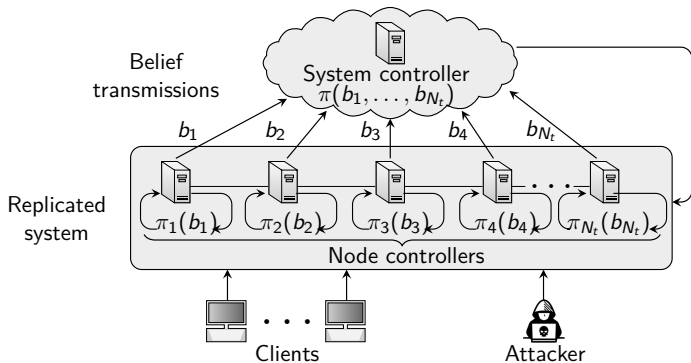
- ▶ The local level models intrusion recovery.
- ▶ The global level models replication control.

Assumption 1

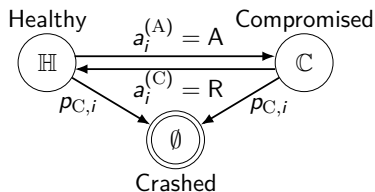
The probability that the system controller fails is negligible.

Assumption 2

Compromise and crash events are statistically independent across nodes.



The Local Control Problem



- ▶ **Partially observed Markov decision process** Γ_i .
- ▶ **Controller actions:** (R)ecover and (W)ait. $a_{i,t} \in \{R, W\}$.
- ▶ **Node states:** $\mathcal{S}_N = \{(\text{H})\text{ealthy}, (\text{C})\text{ompromised}, \emptyset\}$. $s_{i,t} \in \mathcal{S}_N$.
- ▶ **State transition function:** $f(s_{i,t} | s_{i,t}, a_{i,t})$.
- ▶ $p_{C,i}$: crash probability, $p_{A,i}$: intrusion probability.
- ▶ **Observation** $o_{i,t} \sim z_i(\cdot | s_{i,t})$: e.g., IDS alerts at time t .

Node Controller Strategy

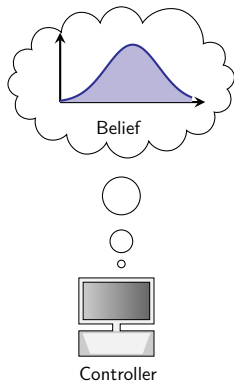
- ▶ The controller computes the **belief**

$$b_{i,t}(s) \triangleq \mathbb{P}[S_{i,t} = s | \mathbf{h}_t].$$

$$\mathbf{h}_t \triangleq (b_{i,1}, a_{i,1}, o_{i,2}, a_{i,2}, o_{i,3}, \dots, a_{i,t-1}, o_{i,t}).$$

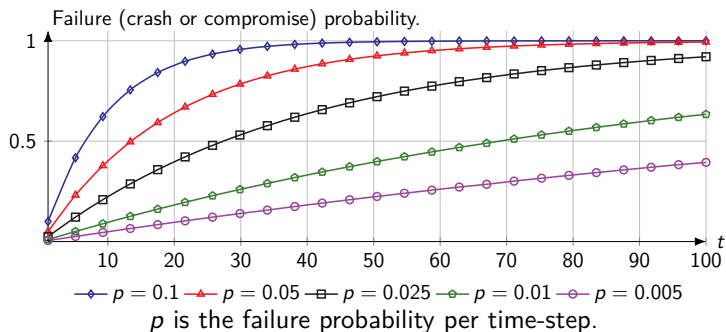
- ▶ **Controller strategy:**

$$\pi : [0, 1] \rightarrow \{W, R\}.$$

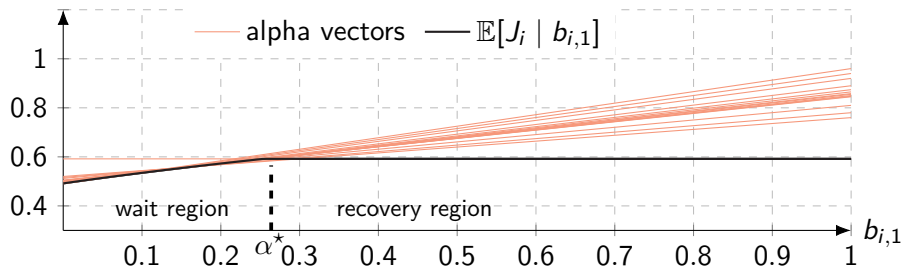


Node Controller Objective

- ▶ **Cost:** $J_i \triangleq \eta T_i^{(R)} + F_i^{(R)}$.
 - ▶ $T_i^{(R)}$ is the average *time-to-recovery*.
 - ▶ $F_i^{(R)}$ is the *recovery frequency*.
 - ▶ $\eta > 1$ is a scaling factor.
- ▶ **Bounded-time-to-recovery constraint:** The time between two recoveries can be at most Δ_R .



Threshold Structure of the Optimal Control Strategy



The controller's optimal cost function.

Theorem 2

There exists an optimal control strategy that satisfies

$$\pi_{i,t}^*(b_{i,t}) = R \iff b_{i,t} \geq \alpha_{i,t}^* \quad \forall t,$$

where $\alpha_{i,t}^ \in [0, 1]$ is a threshold.*

Efficient Computation of Optimal Recovery Strategies

Algorithm 1: Threshold Optimization

1 **Input:** Objective function J_i , parametric optimizer po .

2 **Output:** An approximate optimal control strategy $\hat{\pi}_{i,\theta}$.

3 Algorithm

4 $\Theta \leftarrow [0, 1]$.

5 For each $\theta \in \Theta$, define $\pi_{i,\theta}(b_{i,t})$ as

6
$$\pi_{i,\theta}(b_{i,t}) \triangleq \begin{cases} R & \text{if } b_{i,t} \geq \theta \\ W & \text{otherwise.} \end{cases}$$

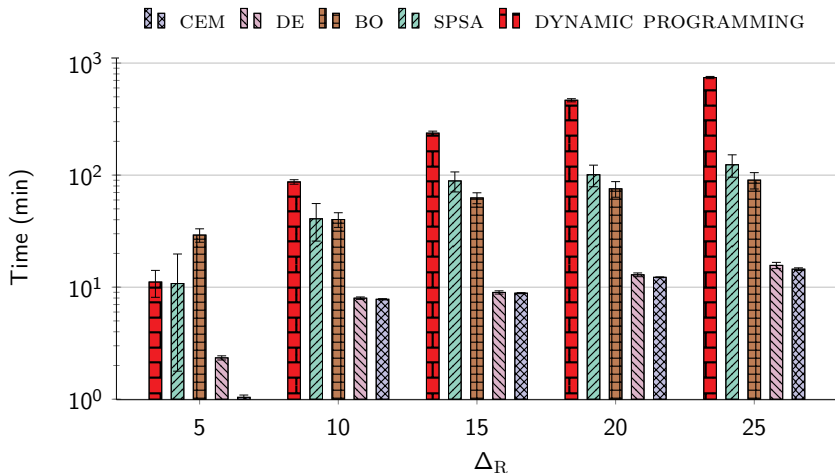
7 $J_\theta \leftarrow \mathbb{E}_{\pi_{i,\theta}}[J_i]$.

8 $\hat{\pi}_{i,\theta} \leftarrow \text{po}(\Theta, J_\theta)$.

9 **return** $\hat{\pi}_{i,\theta}$.

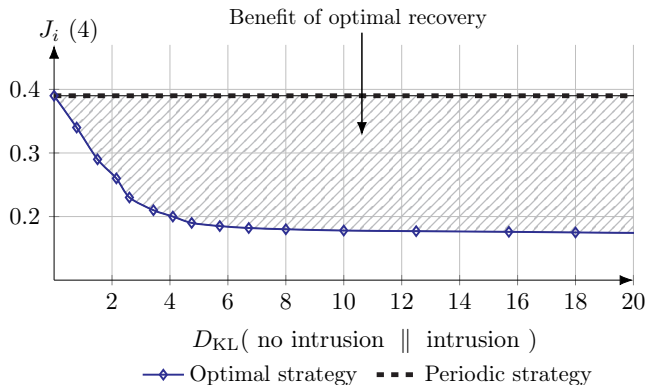
- ▶ Examples of **parameteric optimization algorithmns**: CEM, BO, CMA-ES, DE, SPSA, etc.

Efficient Computation of Optimal Recovery Strategies



Mean compute time to obtain an optimal recovery strategy for different values of the bounded-time-to-recovery constraint Δ_R .

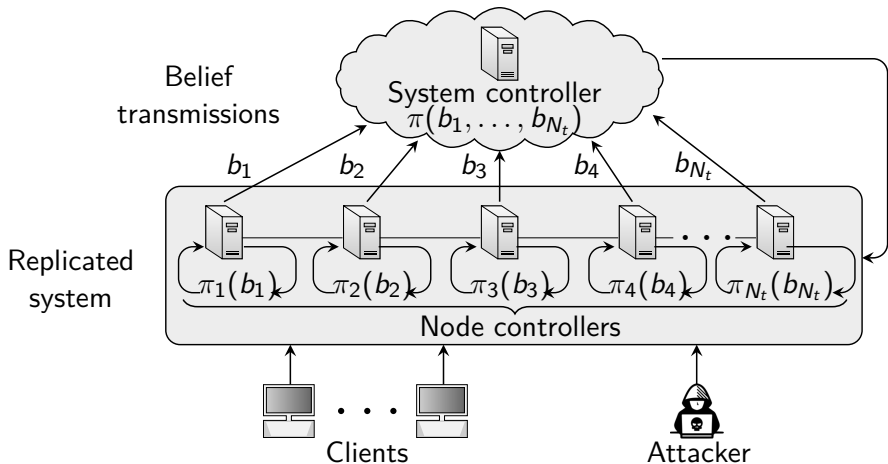
The Benefit of Optimal Recovery Control



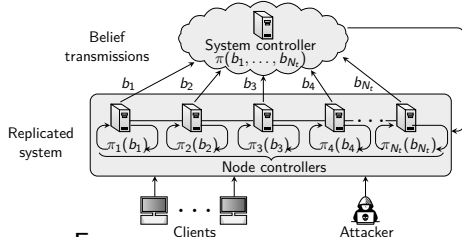
Key insight

Optimal recovery control **can significantly reduce operational cost** given that an **intrusion detection model is available**.

Intrusion Tolerance as a Two-Level Control Problem



The Global Control Problem



- ▶ **Constrained Markov decision process** Γ .
- ▶ **States:** $\mathcal{S}_S = \{0, 1, \dots, s_{\max}\}$, the number of healthy nodes.
- ▶ **Controller actions:** Add $a_t^{(C)} \in \{0, 1\}$ nodes.
- ▶ **Dynamics** f : depend on the local nodes.
- ▶ **Markov** strategy:

$$\pi : \mathcal{S}_S \rightarrow \{0, 1\}.$$

System Controller Objective

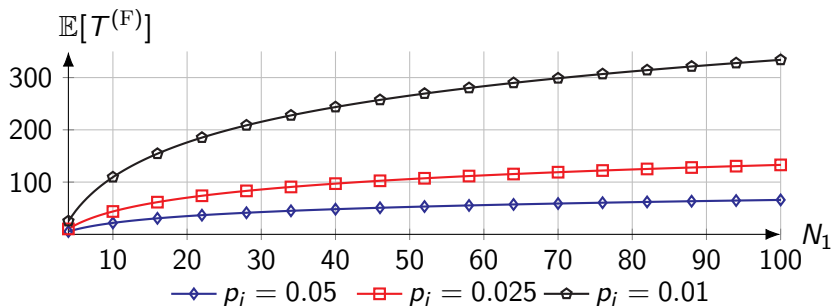
- ▶ Cost: $J \triangleq \lim_{T \rightarrow \infty} \sum_{t=1}^T \frac{a_t}{T}$.
- ▶ Constraint: $T^{(A)} \geq \epsilon_A$, where $T^{(A)}$ is the availability.

ϵ_A	<i>Allowed service downtime per year</i>
0.9	36 days
0.95	18 days
0.99	3 days
0.999	8 hours
0.9999	52 minutes
0.99999	5 minutes
1	0 minutes

System Reliability Analysis

- ▶ The **Mean-time-to-failure** (MTTF) is the **mean hitting time** of a state where $s_t \leq f$:

$$\mathbb{E}[T^{(F)} \mid S_1 = s_1] = \mathbb{E}_{(S_t)_{t \geq 1}} \left[\inf \{t \geq 1 \mid S_t \leq f\} \mid S_1 = s_1 \right].$$



The MTTF in function of the number of initial nodes N_1 and failure probability per node p_i .

Theorem 3 (Optimal Control Strategy Existence)

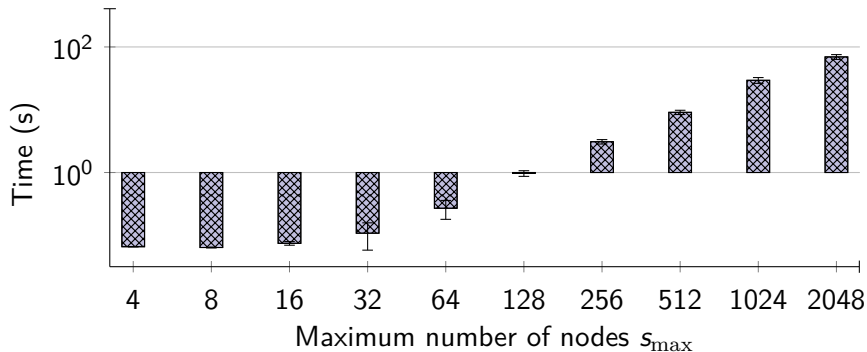
Assuming

- (A) The Markov chain induced by any control strategy is **unichain**.
- (B) The *availability constraint is feasible*.

Then the following holds.

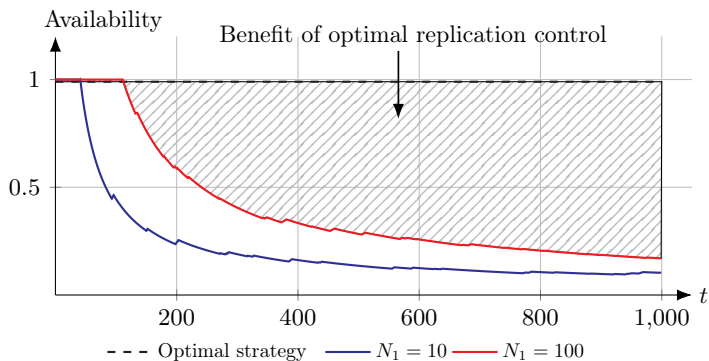
1. There *exists an optimal stationary replication control strategy*.
2. The optimal strategy has a *threshold structure*.
3. An optimal replication control strategy can be computed by using **linear programming**.

Efficient Computation of Optimal Replication Control Strategies



Mean compute time to obtain an optimal replication control strategy.

The Benefit of Optimal Replication Control

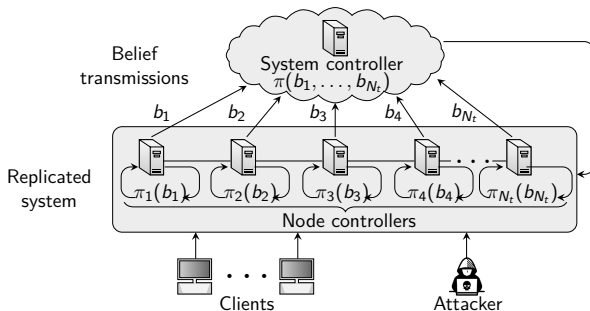


Key insight

Optimal replication control can **guarantee a high service availability in expectation**. The **benefit of optimal replication is mainly prominent for long-running systems**.

Summary of the Control-Theoretic Model

- ▶ **Intrusion recovery control.**
 - ▶ **Partially observed Markov decision process.**
 - ▶ **Threshold structure** of optimal control strategies.
 - ▶ Efficient computation through *stochastic approximation*.
- ▶ **Replication control.**
 - ▶ **Constrained Markov decision process.**
 - ▶ **Threshold structure** of optimal control strategies.
 - ▶ Efficient computation through *linear programming*.

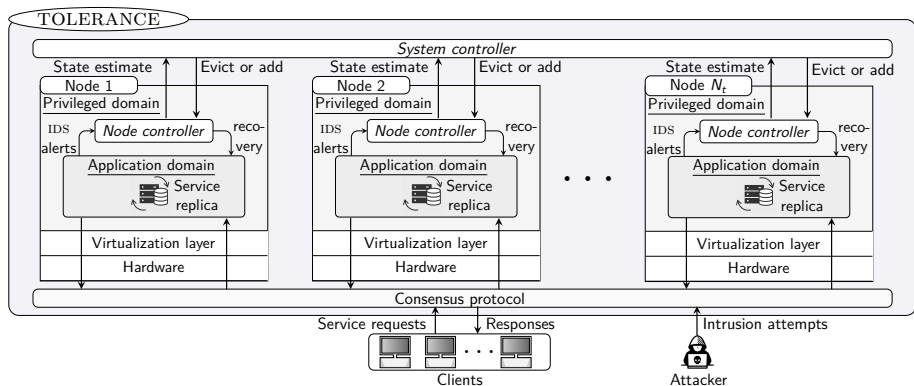


Experiment Setup - Testbed



The TOLERANCE Architecture

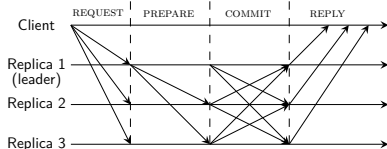
Two-level recovery and replication control with feedback.



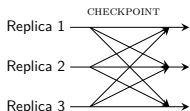
- ▶ A replicated **web service** which offers two operations:
 - ▶ A **read** operation that returns the service state.
 - ▶ A **write** operation that updates the state.

Intrusion-Tolerant Consensus Protocol (MINBFT)

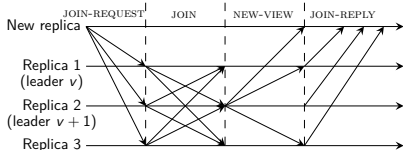
a) Normal operation



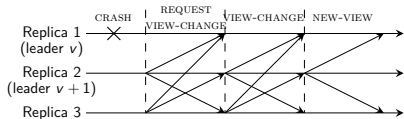
c) Checkpoint



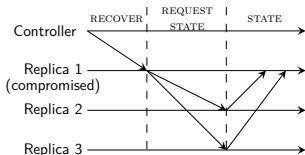
e) Join



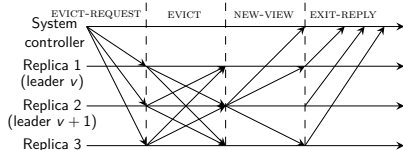
b) View change



d) State transfer



f) Evict



Experiment Setup - Emulated Intrusions

<i>Replica ID</i>	<i>Intrusion steps</i>
1	TCP SYN scan, FTP brute force
2	TCP SYN scan, SSH brute force
3	TCP SYN scan, TELNET brute force
4	ICMP scan, exploit of CVE-2017-7494
5	ICMP scan, exploit of CVE-2014-6271
6	ICMP scan, exploit of CWE-89 on DVWA
7	ICMP scan, exploit of CVE-2015-3306
8	ICMP scan, exploit of CVE-2016-10033
9	ICMP scan, SSH brute force, exploit of CVE-2010-0426
10	ICMP scan, SSH brute force, exploit of CVE-2015-5602

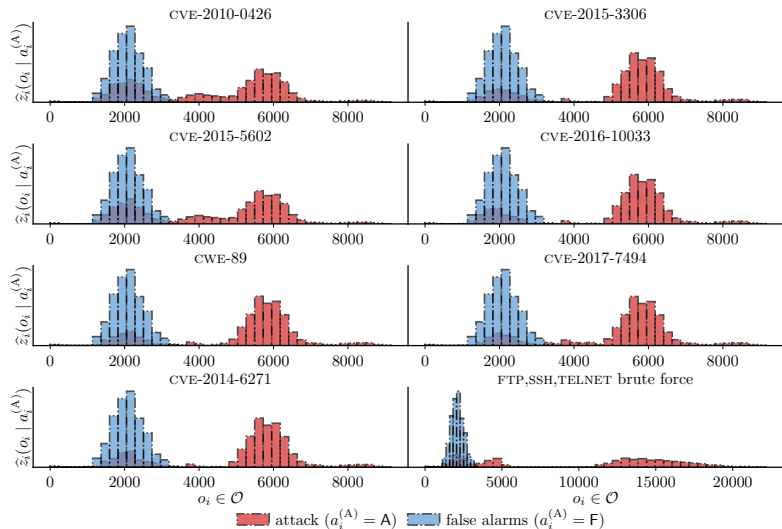
Table 1: Intrusion steps.

Experiment Setup - Background Traffic

<i>Background services</i>	<i>Replica ID(s)</i>
FTP, SSH, MONGODB, HTTP, TEAMSPEAK	1
SSH, DNS, HTTP	2
SSH, TELNET, HTTP	3
SSH, SAMBA, NTP	4
SSH	5, 7, 8, 10
DVWA, IRC, SSH	6
TEAMSPEAK, HTTP, SSH	9

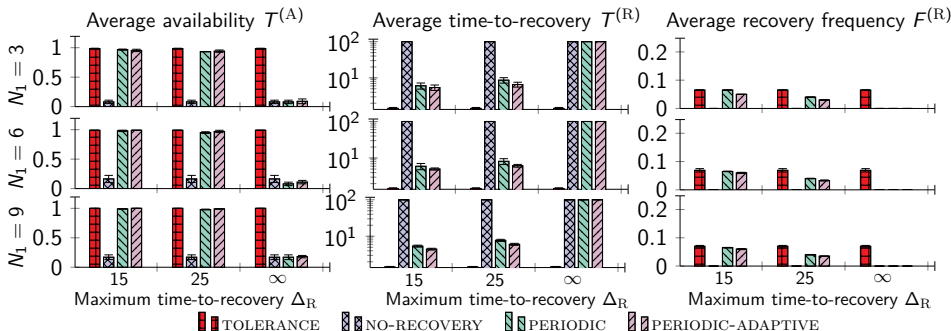
Table 2: Background services.

Estimated Distributions of Intrusion Alerts



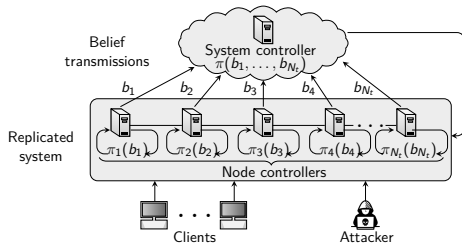
- ▶ We estimate the observation distribution z with the empirical distribution \hat{Z} based on M samples.
- ▶ $\hat{Z} \xrightarrow{a.s.} z$ as $M \rightarrow \infty$ (Glivenko-Cantelli theorem).

Comparison with State-of-the-art Intrusion-Tolerant Systems



Comparison between our optimal control strategies and the baselines; x-axes indicate values of Δ_R ; rows relate to the number of initial nodes N_1 .

Conclusion



- ▶ We present a **control-theoretic model of intrusion tolerance**.
- ▶ We establish **structural results**.
- ▶ We **evaluate the optimal control strategies on a testbed**.
- ▶ Our control-theoretic strategies have **stronger theoretical guarantees and significantly better practical performance** than the heuristic control strategies used in state-of-the-art intrusion-tolerant systems.